# Computing for Genomics 

## Paint by Numbers and Beyond

## Genomics - The study of Genomes

Biological methods
for capture

Analytical methods for interpretation

Recombinant technologies for generation

## The Genome


http://www.healthline.com/health-blogs/tech-medicine/creating-dna-art
http://i.livescience.com/images/i/000/017/621/i02/ITC_EukaryoticCell_Copy.jpg?1309355705

## The Genome

- Genes (coding DNA)
- Noncoding
- Regulatory regions
- Structural
- Repeat elements
- Non coding RNA
- Pseudogenes/Relics/Unclassified


## Repeat Elements

- Genes
- Introns
- Other Intergenic
- DNA Transposon
- Simple Repeats
- Segmental Duplications
- LINEs

SINEs
LTR retrotransposons


Gregory, T. Ryan. "Synergy between sequence and size in large-scale genomics." Nature Reviews Genetics 6.9 (2005): 699-708.

## Human Variation

- Point mutations
- SNPs
- Small insertions, deletions and indels
- Structural Variation
- Copy Number Variation


## The Human Reference Genome

- Around 3Gb
- Haploid ( one version of each chromosome )
- Aims to be point of reference for research
- Publically available
- Consistent coordinates
- Lots of annotation
- Well documented major and minor releases


Morey, Marcos, et al. "A glimpse into past, present, and future DNA sequencing." Molecular genetics and metabolism 110.1 (2013): 3-24.

## Genome Assembly



Lander, Eric S., et al. "Initial sequencing and analysis of the human genome." Nature 409.6822 (2001): 860-921. Venter, J. Craig, et al. "The sequence of the human genome." Science 291.5507 (2001): 1304-1351.

## Sequence Alignment

- Find the best approximate match
- Global

- Local

- Free-end Global

$$
\begin{array}{cccccc}
A & C & A & A & C & G \\
& & \mid & X & \text { I } & \\
- & - & A & G & C & -
\end{array}
$$

## Global Alignment



| match | 1 |
| ---: | :--- |
| mismatch | -3 |
| gap | -2 |\(\quad D(i, j)=\max \left\{\begin{array}{l}D(i-1, j-1)+(match || mismatch) <br>

D(i-1, j)+gap <br>
D(i, j-1)+gap\end{array}\right.\)

Backtrack to get alignment: A C A A C G


## Local Alignment



| match | 3 |
| ---: | :---: |
| mismatch | -2 |
| gap | -3 |\(\quad M(i, j)=\max \left\{\begin{array}{l}0 <br>

M(i-1, j-1)+(match || mismatch) <br>
M(i-1, j)+gap <br>
M(i, j-1)+gap\end{array}\right.\)

Backtrack to get alignment: A C A A C G A G C

## Local Alignment

- Longest common subsequence
- Will find the best aligning substring in both
- i.e. may not align the whole read

```
AGATGTGCTGCCGCC
TTTGTACTGAAA
```


## Free-end Global Alignment

|  |  | A | C | A | A | C | G |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| A | 0 | 1 | -1 | 1 | 1 | -3 | -3 |
| G | 0 | -2 | -2 | -1 | -1 | -2 | -2 |
| C | 0 | -2 | -1 | -3 | -2 | 0 | -2 |


| match | 1 |
| ---: | :--- |
| mismatch | -3 |
| gap | -2 |\(\quad D(i, j)=\max \left\{\begin{array}{l}D(i-1, j-1)+(match || mismatch) <br>

D(i-1, j)+gap <br>
D(i, j-1)+gap\end{array}\right.\)
$\begin{array}{llllllll}\text { Backtrack to get alignment: } & \text { A } & \text { C } & \text { A } & \text { A } & \text { C } & \text { G } \\ & & & \mid & \times & \mid & \\ & - & - & A & G & C & -\end{array}$

## Free-End Alignment

- Aligns whole of both reads
- Containment
- Longest prefix/suffix overlap

TTCAGATGTGCTG
TGTACTGACGTAG

## Dynamic Programming

- $O(\mathrm{~nm})$ for time* and space complexity **


## Genome Assembly



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Venter, J. Craig, et al. "The sequence of the human genome." Science 291.5507 (2001): 1304-1351.

## The Human Reference Genome

- Based on limited subjects
- Does not capture variation
- is one "Golden path"
- Subsequences reported are not unique
- It has all repeats found in these subjects that could be resolved


## Next Generation Sequencing

- Traditional sequencing doesn't scale up
- Next Generation Sequencers
- high throughput (4h-3days)
- high coverage (20x-50x)
- short reads
(25-200bp)

eepto, Itumina inc. Al righto recerved.


## NGS Uses Resequencing

- NGS produces huge amounts of data
- 120Gb-1Tb compressed
- Dynamic Programming is impractical
- Rather than assemble:
- map to the reference quickly
$\rightarrow$ Read mapping algorithms
- verify local alignment (and call variants)
$\rightarrow$ Dynamic programming
$\rightarrow$ Local de-novo assembly


## Exact Matching

- Instead of looking for "good" matches, only look for all exact matches

Still not enough for genome scale

- Exact string matching algorithms time complexity
- Worst: O(nm)
- Best: $\quad \Omega(n / m)$


## Fast Approximate Matching

- Expect very few differences between the sample's reads and the reference genome
- Sequencing errors
- Natural variation
- Expect even fewer differences between sample's reads
- Sequencing error
- May be variation within a sample (tissue)
- May be variation between repeated regions


## Read Mapping Algorithms

Two main approaches

- Filter
- Index


## Filtering

- Reduce the number of possible approximate matches
- In practice, want really good alignments
$\rightarrow$ Expect sections of the alignment to be exact
$\rightarrow$ Expect no more than $k$ errors (mismatches)


## Pigeonhole Lemma



- Assume no more than k errors tolerated
- Divide query into k+1 pieces
- Search for each of these in the genome
- If the query is in the genome, one will match exactly
- Report that as a candidate region


## q-gram Lemma



- Assume no more than $k$ errors tolerated
- Create all possible overlapping q-grams from the read
- search for all of these


## q-gram Lemma



- Number of q-grams for read of length n ?
- k errors affect how many q-grams at most in worst case?
- Assume no more than $k$ errors tolerated
- Create all possible overlapping q-grams from the read
- search for all of these


## q-gram Lemma



- Number of q-grams for read of length n ?
- k errors affect how many q-grams at most in worst case?
- Assume no more than k errors tolerated
- Create all possible overlapping q-grams from the read
- search for all of these
- If the query is in the genome, at least $n-(k+1) q+1$ of the $q$-grams match exactly
- count the number of q-grams that matched, and if it passes this threshold, report a candidate region


## Finding Candidate Regions



## Indexing

- Store the genome/reads in a data structure that facilitates fast exact or near-exact alignment
- Must be reasonable for memory limits of the machine

Genome(s) (cont.)
overview of, 24-25, 25 f
plasmid, 26, 37, 37f, 41f prokaryotic, 38-39, 38f repeated sequences in, 288-294, 301-303
sequencing of. See Genomic sequencing
size of, 35, 36f, 36t, 266, 266f, 266t, 306, 386 evolution and, 374
structure of, 301-303
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transposable elements in, 249, 249f, 250, 289-290, 290f, 302
viral, 26, 37f, 38, 41f
Genome projects, 268, 270-271, 271t
databases for, 270-271, 271t, 651-653
Genomic clones
in contig, 281, 281f
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by fluorescent in situ hybridization, 284, 285f
databases for, 651-653
filling gaps in, 294
genomic subdivision for, 283
goals of, 286-287
information gaps regarding, 295-296, 296f
interpretation problems in, 295-296, 296f
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mobile genetic elements and, 289-291, 290f, 291f
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paired-end sequences in, 287f, 293
prediction of mRNA and polypeptide structure from, 296-301
primer walking in, 287
purposes of, 270-271
scaffolds in, 292f, 293
steps in, 292 f
tandem array repeats and,
288-289, 288f, 289f
whole genome shotgun,
267, 287, 292f,
293-294

## Indexing Data Structures


\$ B A N A N A
A\$BANAN ANA\$BAN ANANA\$B B A N A N A \$ NA\$BANA NANA\$BA

## FM Index

- Uses Burrows-Wheeler transform
- Plus extra tables to speed things up



## Burrows-Wheeler transform

- Because they are rotations
- the character in the last column is what precedes the character in the first column in the original string
- Because the suffixes are sorted
- the first row, last character is the end of the original string



## Burrows-Wheeler transform

- Because the suffixes are sorted (cont'd)
- the rank of the character in the last column is the same as the first

\$acaac $g_{0}$ $\mathrm{a}_{0} \mathrm{acg}_{\mathrm{c}} \mathrm{ac}_{0}$<br>a,caacg \$<br>$\mathrm{a}_{2} \mathrm{c}$ g \$ acta<br>$c_{0}$ atcg \$ $\mathrm{a}_{1}$<br>$c_{1}$ g \$aca $a_{2}$<br>$\mathrm{g}_{0}$ \$acaac

## Burrows-Wheeler transform

- We can find the position in the first column (F) based only on information from the last column (L)

$$
\begin{aligned}
\mathrm{LF}\left({ }^{\prime} c^{\prime}, 6\right) & =\operatorname{Occ}\left({ }^{\prime} c^{\prime}\right)+\operatorname{Count}\left({ }^{\prime} c^{\prime}, 6\right) \\
& =4+1
\end{aligned}
$$

\$acaacgo $\mathrm{a}_{0} \mathrm{acg}_{\mathrm{c}} \mathrm{ac}_{0}$ $a_{1} c$ acg \$ $a_{2} \mathbf{c}$ g \$ acta $c_{0}$ a acg \$ $\mathrm{a}_{1}$ $c_{1}$ g \$aca $a_{2}$ $\mathrm{g}_{0} \$$ acaac ${ }_{1} \downarrow$

## Burrows-Wheeler transform

- Walk left algorithm
- We can use the BWT and the LF function to reconstruct the original text

$$
\begin{aligned}
& i=0 \\
& t=">
\end{aligned}
$$

while bwt[i] != ‘\$':

$$
\begin{aligned}
& \mathbf{t}=\operatorname{bwt}[i]+\mathbf{t} \\
& \mathbf{i}=\operatorname{LF}(\mathbf{i}, \operatorname{bwt}[i])
\end{aligned}
$$

| g | c g | acg | a acg | caacg | $\overbrace{\text { acaacg }}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| \$ acaacg | \$ acaacg | \$ acaacg | \$ a c a acg | \$ acaacg | \$ acaacg |
| a acg ${ }^{\text {a }}$ c | a acg \$/c | a acg ${ }^{\text {a }}$ c | a a c g \$ ac | a -y ¢ | a a c g \$ zc |
| a caacg \$ | aca a g \$ | acaacg \$ | acaacg \$ | a ca~g \$ | acaalg \$ |
| a c g \$ a c a | $\mathbf{a} c \mathrm{~g} / \mathrm{ac} \mathbf{a}$ | a c g \$ a c a | $a \rightarrow a$ | a c g \$ a ca | ac.s a c a |
| caacg\$ ${ }^{\text {a }}$ | c a 0 | caacg \$ $\mathbf{a}$ |  | ca acg \$ a | caucy ${ }^{\text {a }}$ |
| g \$ a c a $a$ | c $c$ a c a $\mathbf{a}$ | c | c g \$ aca | c g \$ a c a $\mathbf{a}$ | c g \$ a c a $\mathbf{a}$ |
| g \$ a caac | $\mathrm{g} \rightarrow$ ¢0a | g \$ a caac | g\$aca ac | g\$acaac | g \$ a caa |

## FM Index for Exact Matching



## FM Index for Exact Matching



$$
\begin{aligned}
& q=\text { "aac" } \\
& \text { top }=0 \\
& \text { bot = len(bwt) } \\
& \text { for qc in reverse(q): } \\
& \text { top }=\operatorname{LF}(\text { top, qc) } \\
& \text { bot }=\operatorname{LF}(b o t, \\
& q c)
\end{aligned}
$$

## FM Index for Exact Matching

| aac | aac | a ac | a ac |
| :---: | :---: | :---: | :---: |
| $\rightarrow$ acaacg | \$ acaacg | \$ acaacg | \$ acaacg |
| a acg \$ a c | a acg \$ac | a acg \$ac | a acg $\mathrm{ac}_{\text {c }}$ |
| acaacg\$ | acaacg \$ | acaacg\$ | acaacg\$ |
| $\mathbf{a c g \$ a c} \mathbf{a}$ | $\mathbf{a c g \$ a c a}$ | acg\$aca | $\mathbf{a} \mathrm{cg} \$ \mathrm{ac} \mathbf{a}$ |
| caacg \$a | caacg \$ | caacg\$a | caacg \$a |
| c g \$acaa | cg\$acaa | c g \$acaa | c g \$ acau |
|  |  | g\$acaac | g\$acaac |

Track a top and bottom index

- These bound the remaining possible matches at each step
- If they are ever the same, there are no matches

Here, LF is called on the query characters

- "Where are the first \& last qc you saw in F?"
$\mathrm{q}=$ "aac"
top $=0$
bot $=\operatorname{len}(b w t)$
for qc in reverse(q):
top $=L F($ top, qc)
bot $=\operatorname{LF}($ bot, qc)

The end TOP index is our match!

## FM Index for Exact Matching



Track a top and bottom index

- These bound the remaining possible matches at each step
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Here, LF is called on the query characters

- "Where are the first \& last qc you saw in F?"

The end TOP index is our match!

## FM Index

- How do we find this in the genome?
- Isn't the count operation $O(n)$ ?


## FM Index

- How do we find this in the genome?

1) Can use our walk left algorithm to reconstruct

| 6 |  |
| :---: | :---: |
| 2 | acg ${ }^{\text {a }}$ |
| 0 | caacg\$ |
| 3 | c g \$aca |
| 1 | caacg\$ |
| 4 |  |
| 5 | a |

## FM Index

- How do we find this in the genome?

1) Can use our walk left algorithm to reconstruct

| 6 | \$acaacg |
| :---: | :---: |
| 2 | aacg\$ac |
| 0 | caacg\$ |
| 3 | c g \$ ac |
| 1 | caacg\$ |
| 4 |  |
| 5 | g\$acaac |

2) Could store the entire suffix array

## FM Index

- How do we find this in the genome?

1) Can use our walk left algorithm to reconstruct

| 6 |  |
| :---: | :---: |
| 2 | a acg\$ac |
| 0 | acaacg |
| 3 | acg\$ac |
| 1 | caacg\$ |
| 4 | a |
| 5 | g\$ acaac |

2) Could store the entire suffix array
3) Only store certain rows of (2), use (1) until we get to one.

## FM Index

- Isn't the count operation $O(n)$ ?


## FM Index

- Isn't the count operation $O(n)$ ?
- We again store cumulative counts for certain rows, for each of \$ACTG
- Also, the reference is usually put in backward (or both forward and backward)


## Beyond?

- Read-mapping is limited by
- the reference genome assembly
- exact matching
- So, why not assemble each time?


## De Novo Assembly

- Overlap-Layout-Consensus
- De Bruijn Graphs


## Overlap-Layout-Consensus

- Overlap
- Compute overlap score of all reads
- Layout
- Create a graph where nodes are a read, edges are overlaps between reads
- Find their "layout" by finding a Hamiltonian path through the graph
- Consensus
- Find the consensus sequence by reading nodes along the path



## De Bruijn Graphs

- Break reads into k-mers
- Each node in the graph is a k-mer
- Connect an edge to the next k-mer found in the read
- Find a Eulerian path


## De Novo Assembly

- Confounded by repeats
- repeats are collapsed in these representations
- may have many ways in and out of these graph regions.


## Further beyond

- Newer sequencing techniques
- Nanopore sequencing
- Single cell sequencing
- Downstream analysis issues
- How do we compare genomes?
- How do we store them?
- Improve the reference model


## Further beyond

- Downstream analysis issues
- How do we compare genomes?
- How do we store them?

https://en.wikipedia.org/wiki/1000_Genomes_Project\#/media/ File:Genetic_Variation.jpg
- Improve the reference model


## Thank you!

References and resources

- Algorithms
- Reinert, Knut, et al. "Alignment of Next-Generation Sequencing Reads." Annual review of genomics and human genetics 0 (2015).
- Li, Heng, and Nils Homer. "A survey of sequence alignment algorithms for next-generation sequencing." Briefings in bioinformatics 11.5 (2010): 473-483.
- Sequencing
- Morey, Marcos, et al. "A glimpse into past, present, and future DNA sequencing." Molecular genetics and metabolism 110.1 (2013): 3-24.
- Ben Langmead's Teaching Resources:
- http://www.langmead-lab.org/teaching-materials/

Supervisor: Mark Daley
DaleyLab.org

## Structural Variant Discovery



Split reads


