Computing for Genomics

Paint by Numbers and Beyond

Genomics – The study of Genomes

Biological methods for capture

Analytical methods for interpretation

Genomics

Recombinant technologies for generation

The Genome



- The DNA in each cell
- ~3 billion base pairs
- Any two people are 99.6% to 99.9% the same



http://www.healthline.com/health-blogs/tech-medicine/creating-dna-art http://i.livescience.com/images/i/000/017/621/i02/ITC_EukaryoticCell_Copy.jpg?1309355705

The Genome

- Genes (coding DNA)
- Noncoding
 - Regulatory regions
 - Structural
 - Repeat elements
 - Non coding RNA
 - Pseudogenes/Relics/Unclassified

Repeat Elements



- Introns
- Other Intergenic
- DNA Transposon
- Simple Repeats
- Segmental Duplications
- LINEs
- SINEs
- LTR retrotransposons



Gregory, T. Ryan. "Synergy between sequence and size in large-scale genomics." Nature Reviews Genetics 6.9 (2005): 699-708.

Human Variation

- Point mutations
 - SNPs
- Small insertions, deletions and indels
- Structural Variation
 - Copy Number Variation

The Human Reference Genome

• Around 3Gb

- Haploid (one version of each chromosome)

- Aims to be point of reference for research
 - Publically available
 - Consistent coordinates
 - Lots of annotation
 - Well documented major and minor releases



Morey, Marcos, et al. "A glimpse into past, present, and future DNA sequencing." Molecular genetics and metabolism 110.1 (2013): 3-24.

Genome Assembly



Lander, Eric S., et al. "Initial sequencing and analysis of the human genome." Nature 409.6822 (2001): 860-921. Venter, J. Craig, et al. "The sequence of the human genome." Science 291.5507 (2001): 1304-1351.

Sequence Alignment

- Find the best approximate match
 - Global

Local

A C A A C G | X | A - - G C -A C A A C G | X | A C A C G A C C

A C A A C G | X | - - A G C -

- Free-end Global

Global Alignment

	А	С	А	А	С	G
0	-2	-4	-6	-8	-10	-12
-2	1	-1	-3	-5	-7	-9
-4	-3	-2	-4	-6	-8	-6
-6	-5	-2	-4	-6	-5	-7

Α

G

С

match1D(i-1, j-1) + (match || mismatch)mismatch-3D(i,j) = maxD(i-1, j) + gapgap-2D(i, j-1) + gap

Backtrack to get alignment: A C A A C G | X | A - - G C -

Local Alignment

	А	С	А	А	С	G
0	0	0	0	0	0	0
0	3	0	3	3	1	0
0	0	1	0	1	1	4
0	0	3	0	0	4	0

match 3 mismatch -2 gap -3 M(i,j) = max $\begin{bmatrix} 0\\ M(i-1, j-1) + (match || mismatch)\\ M(i-1, j) + gap\\ M(i, j-1) + gap\end{bmatrix}$ Backtrack to get alignment: A C A A C G

A G C

Local Alignment

- Longest common subsequence
- Will find the best aligning substring in both
 i.e. may not align the whole read

AGATGTGCTGCCGCC |||×||| TTTGTACTGAAA

Free-end Global Alignment

	А	С	А	А	С	G
0	0	0	0	0	0	0
0	1	-1	1	1	-3	-3
0	-2	-2	-1	-1	-2	-2
0	-2	-1	-3	-2	0	-2

Α

G

С

match1D(i-1, j-1) + (match || mismatch)mismatch-3D(i,j) = maxD(i-1, j) + gapgap-2D(i, j-1) + gap

Backtrack to get alignment: A C A A C G | X | - - A G C -

Free-End Alignment

- Aligns whole of both reads
 - Containment
 - Longest prefix/suffix overlap



Dynamic Programming

• O(nm) for time* and space complexity **

Genome Assembly



Lander, Eric S., et al. "Initial sequencing and analysis of the human genome." Nature 409.6822 (2001): 860-921. Venter, J. Craig, et al. "The sequence of the human genome." Science 291.5507 (2001): 1304-1351.

The Human Reference Genome

- Based on limited subjects
 - Does not capture variation
 - is one "Golden path"
- Subsequences reported are not unique
 - It has all repeats found in these subjects that could be resolved

Next Generation Sequencing

• Traditional sequencing doesn't scale up

- Next Generation Sequencers
 - high throughput (4h-3days)
 - high coverage (20x-50x)
 - short reads (25-200bp)

Automated Workflow



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NGS Uses Resequencing

- NGS produces huge amounts of data – 120Gb - 1Tb compressed
- Dynamic Programming is impractical
- Rather than assemble:
 - map to the reference quickly
 - → Read mapping algorithms
 - verify local alignment (and call variants)
 - ➔ Dynamic programming
 - →Local de-novo assembly

Exact Matching

 Instead of looking for "good" matches, only look for all exact matches

Still not enough for genome scale

- Exact string matching algorithms time complexity
 - Worst: O(nm)
 - Best: $\Omega(n/m)$

Fast Approximate Matching

- Expect very few differences between the sample's reads and the reference genome
 - Sequencing errors
 - Natural variation
- Expect even fewer differences between sample's reads
 - Sequencing error
 - May be variation within a sample (tissue)
 - May be variation between repeated regions

Read Mapping Algorithms

- Two main approaches
 - Filter
 - Index

Filtering



- Reduce the number of possible approximate matches
 - In practice, want really good alignments
 Expect sections of the alignment to be exact
 Expect no more than k errors (mismatches)

Pigeonhole Lemma



Red shows mismatch to reference (not pictured)

- Assume no more than k errors tolerated
- Divide query into k+1 pieces
 - Search for each of these in the genome
- If the query is in the genome, one will match exactly
 - Report that as a candidate region

q-gram Lemma



- Assume no more than k errors tolerated
- Create all possible overlapping q-grams from the read
 - search for all of these

q-gram Lemma



- Number of q-grams for read of length n?
- k errors affect how many q-grams at most in worst case?

- Assume no more than k errors tolerated
- Create all possible overlapping q-grams from the read
 - search for all of these

q-gram Lemma



- Number of q-grams for read of length n?
- k errors affect how many q-grams at most in worst case?

- Assume no more than k errors tolerated
- Create all possible overlapping q-grams from the read
 - search for all of these
- If the query is in the genome, at least n-(k+1)q +1 of the q-grams match exactly
 - count the number of q-grams that matched, and if it passes this threshold, report a candidate region

Finding Candidate Regions

Pigeonhole

Reference Genome



q-gram



Indexing

- Store the genome/reads in a data structure that facilitates fast exact or near-exact alignment
- Must be reasonable for memory limits of the machine

Genome(s) (cont.) overview of, 24-25, 25f plasmid, 26, 37, 37f, 41f prokaryotic, 38-39, 38f repeated sequences in, 288-294, 301-303 sequencing of. See Genomic sequencing size of, 35, 36f, 36t, 266, 266f, 266t, 306, 386 evolution and, 374 structure of, 301-303 subdivision of, 283-284, 284f transposable elements in, 249, 249f, 250, 289-290, 290f, 302 viral, 26, 37f, 38, 41f Genome projects, 268, 270-271, 271t databases for, 270-271, 271t, 651-653 Genomic clones in contig, 281, 281f ordering of by clone fingerprints, 281, 282f by fluorescent in situ hybridization, 284, 285f

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Indexing Data Structures



Ben Langmead teaching materials: http://www.langmead-lab.org/teaching-materials/

Uses Burrows-Wheeler transform
 – Plus extra tables to speed things up



- Because they are rotations
 - the character in the last column is what precedes the character in the first column in the original string
- Because the suffixes are sorted
 - the first row, last character is the end of the original string

```
$ a c a a c g
                                       $ a c a a c g
                   aacg$ac
                                       a a c g $ a c
                                       acaacg$
                   a c a a c g $
                                  —► acg$ac<mark>a</mark> —► gc$aaac
             — → a c g $ a c a
acaacg$
                                       caacg$a
                   caacg$a
     т
                                                               BWT(T)
                   <mark>c g $</mark> a c a a
                                       cg$aca<mark>a</mark>
                                       g$acaac
                   q S a c a a c
                                         Last column
                       Burrows
                       Wheeler
                        Matrix
```

- Because the suffixes are sorted (cont'd)

 the rank of the character in the last column is the same as the first

```
$ a c a a c g<sub>0</sub>
a<sub>0</sub> a c g $ a c<sub>0</sub>
a<sub>1</sub> c a a c g $
a<sub>2</sub> c g $ a c a<sub>0</sub>
c<sub>0</sub> a a c g $ a<sub>1</sub>
c<sub>1</sub>g $ a c a a<sub>2</sub>
g<sub>0</sub>$ a c a a c<sub>1</sub>
```

• We can find the position in the first column (F) based only on information from the last column (L)

$$LF('c', 6) = Occ('c') + Count('c', 6)$$

= 4 + 1

Occurrence:

Number of letters before any 'c' in F?

4 (\$ and 3 a's)

Count:

How many 'c's have we seen in L?

 $1 (c_0)$

\$ a c a a c g, a,a c g \$ a c, a,c a a c g \$ a,c g \$ a c a, a,c g \$ a c a, c,a a c g \$ a, c,g \$ a c a a, g,\$ a c a a c,

- Walk left algorithm
 - We can use the BWT and the LF function to reconstruct the original text

i = 0
t = ""
while bwt[i] != `\$':
 t = bwt[i] + t
 i = LF(i, bwt[i])

						\frown
	g	c g	a c g	a a c g	c a a c g	a c a a c g
\$	a c a a c <mark>g</mark>	\$ a c a a c g	\$ a c a a c g	\$ a c a a c g	\$ a c a a c g	\$ a c a a c g
а	a c g \$ a c	aacg\$/c	a a c g \$ a c	a a c g \$ a c	a 🔍 e g 🕪 c	aacg\$a_c
а	c a a c g \$	acaagg\$	a caacg \$	acaacg\$	acaacg\$	acaaeg\$
а	c g \$ a c a	acg 🌮 aca	acg\$aca	a ⊲g \$ a ⊳ a	a cg\$a ca	а с 🖋 💲 а с а
С	a a c g \$ a	c a a c g \$ a	c a a c g \$ a	caacg\$a	c a a c g \$ a	c a a c g a
С	g \$ a c a a	c	c g 🗘 a c la a	c g\$aca a	c g\$aca a	c g\$aca a
g	\$ a c a a c	g /⊕ a o a⊯ c	g \$ a c a a c	g \$ a c a a c	g \$ a c a a c	g \$ a c a a c

Final t



Ferragina, Paolo and Manzini, Giovanni. "Opportunistic data structures with applications." Foundations of Computer Science, 2000. Proceedings. 41st Annual Symposium on. IEEE (2000)



Ferragina, Paolo and Manzini, Giovanni. "Opportunistic data structures with applications." Foundations of Computer Science, 2000. Proceedings. 41st Annual Symposium on. IEEE (2000)



Track a top and bottom index

- These bound the remaining possible matches at each step
- If they are ever the same, there are no matches
- Here, LF is called on the *query* characters
 - "Where are the first & last qc you saw in F?"

The end TOP index is our match!

Ferragina, Paolo and Manzini, Giovanni. "Opportunistic data structures with applications." Foundations of Computer Science, 2000. Proceedings. 41st Annual Symposium on. IEEE (2000)

q = "aac" top = 0 bot = len(bwt)for qc in reverse(q): top = LF(top, qc)bot = LF(bot, qc)



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q = "aac" top = 0 bot = len(bwt)for qc in reverse(q): top = LF(top, qc)bot = LF(bot, qc)

• How do we find this in the genome?

• Isn't the count operation O(n)?

• How do we find this in the genome?

1) Can use our walk left algorithm to reconstruct

6	\$ a c a a c <mark>g</mark>
2	a a c g \$ a <mark>c</mark>
0	a c a a c g \$
3	a c g \$ a c <mark>a</mark>
1	c a a c g \$ <mark>a</mark>
4	c g \$ a c a <mark>a</mark>
5	g \$ a c a a <mark>c</mark>

• How do we find this in the genome?

- 1) Can use our walk left algorithm to reconstruct
- 2) Could store the entire suffix array

6	\$ a c a a c <mark>g</mark>
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0	a c a a c g \$
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4	c g \$ a c a <mark>a</mark>
5	g \$ a c a a <mark>c</mark>

- How do we find this in the genome?
- 1) Can use our walk left algorithm to reconstruct
- 2) Could store the entire suffix array
- Only store certain rows of (2), use (1) until we get to one.

	6	\$	а	с	а	а	с	g
	2	а	а	С	g	\$	а	С
	0	а	С	а	а	С	g	\$
Ľ	3	а	С	g	\$	а	С	а
Ŀ	1	С	а	а	С	g	\$	а
4	4	С	g	\$	а	С	а	а
ļ	5	g	\$	а	С	а	а	С

• Isn't the count operation O(n)?

- Isn't the count operation O(n)?
 - We again store cumulative counts for certain rows, for each of \$ACTG

 Also, the reference is usually put in backward (or both forward and backward)

Beyond?

- Read-mapping is limited by
 - the reference genome assembly
 - exact matching
- So, why not assemble each time?

De Novo Assembly

- Overlap-Layout-Consensus
- De Bruijn Graphs

Overlap-Layout-Consensus

Overlap

- Compute overlap score of all reads

- Layout
 - Create a graph where nodes are a read, edges are overlaps between reads
 - Find their "layout" by finding a Hamiltonian path through the graph
- Consensus
 - Find the consensus sequence by reading nodes along the path



De Bruijn Graphs

- Break reads into k-mers
- Each node in the graph is a k-mer
- Connect an edge to the next k-mer found in the read
- Find a Eulerian path

De Novo Assembly

- Confounded by repeats
 - repeats are collapsed in these representations
 - may have many ways in and out of these graph regions.

Further beyond

- Newer sequencing techniques
 - Nanopore sequencing
 - Single cell sequencing
- Downstream analysis issues
 - How do we compare genomes?
 - How do we store them?
- Improve the reference model

Further beyond

- Downstream analysis issues
 - How do we compare genomes?
 - How do we store them?



 Improve the reference model https://en.wikipedia.org/wiki/1000_Genomes_Project#/media/ File:Genetic_Variation.jpg

Thank you!

References and resources

- Algorithms
 - Reinert, Knut, et al. "Alignment of Next-Generation Sequencing Reads." Annual review of genomics and human genetics 0 (2015).
 - Li, Heng, and Nils Homer. "A survey of sequence alignment algorithms for next-generation sequencing." Briefings in bioinformatics 11.5 (2010): 473-483.
- Sequencing
 - Morey, Marcos, et al. "A glimpse into past, present, and future DNA sequencing." Molecular genetics and metabolism 110.1 (2013): 3-24.
- Ben Langmead's Teaching Resources:
 - http://www.langmead-lab.org/teaching-materials/

Supervisor: Mark Daley DaleyLab.org

Structural Variant Discovery

